

07 - Notation-State-Based

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What are we trying to represent?

What?

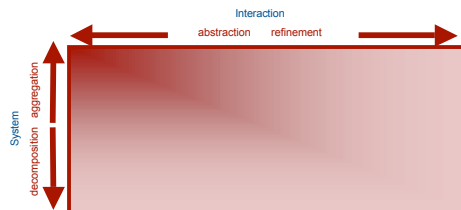
- Decomposition
- Communications
- Functions
- Behavior

Where?

- System level
- Component level

Why?

- No one notation good for representing each
- Need to make connections, make consistent, complete

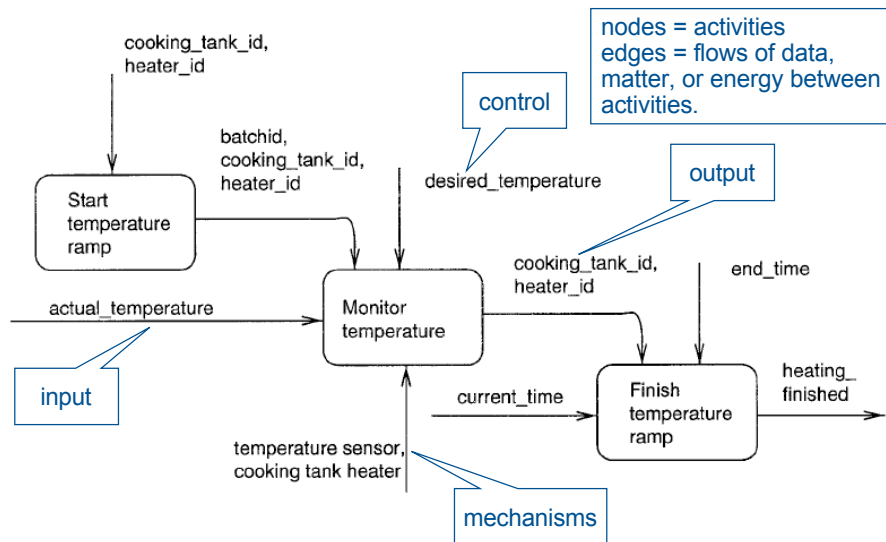


COMPUTER SCIENCE communication specification techniques

- Context Diagrams
- SADT Activity Diagrams
- State Activity Charts
- Object Communication Diagrams
- JSD System Network Diagrams
- SDL Block Diagrams
- Sequence Diagrams
- Collaboration Diagrams

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COMPUTER SCIENCE SADT Activity Diagrams

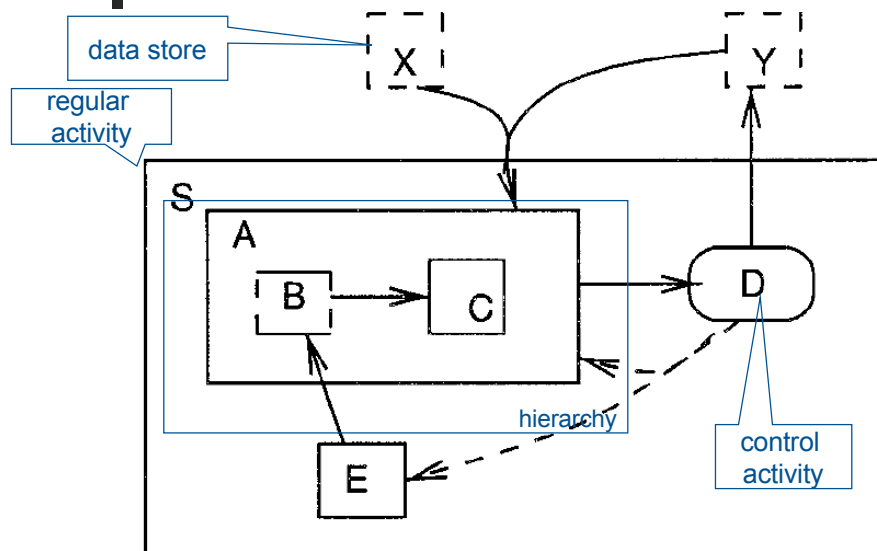


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SADT Activity Diagrams

- well suited to the representation of functional decomposition
 - can be structured hierarchically
 - one activity box into a lower-level activity diagram
 - not a physical but a conceptual decomposition
- In contrast to DFDs
 - activity diagrams have only one kind of component
 - the interfaces between the activities distinguish input, output, control, and mechanism.

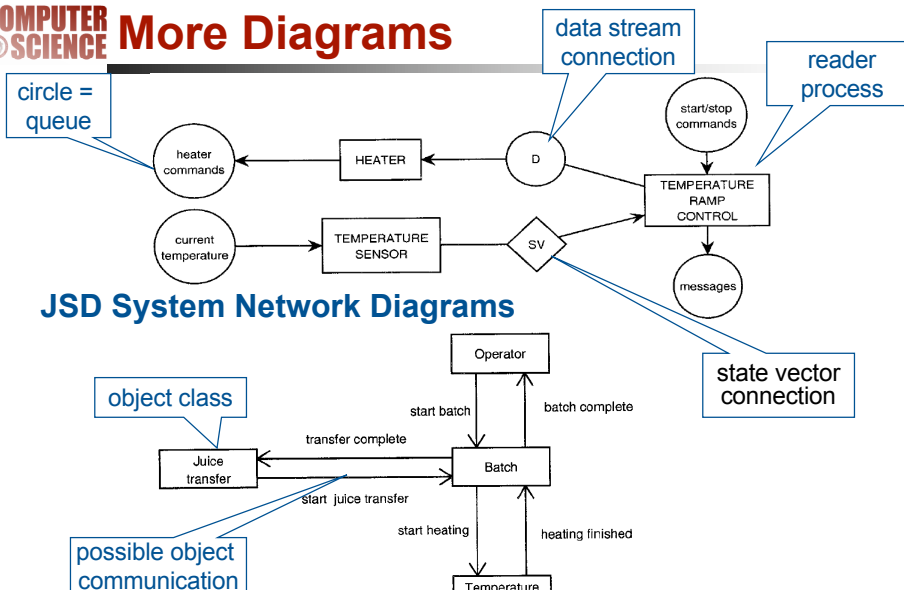
Statemate Activity Charts



State Activity Charts

- Control activities are always sub-activities of regular activities
 - a regular activity can have at most one control activity as an immediate component.
- Control activities are specified by extended finite state machines, called statecharts
 - are finite state machines extended with variables that can be updated.
 - Control activities can thus maintain a state in their variables, can perform data processing (testing and updating the variables), and can contain control.

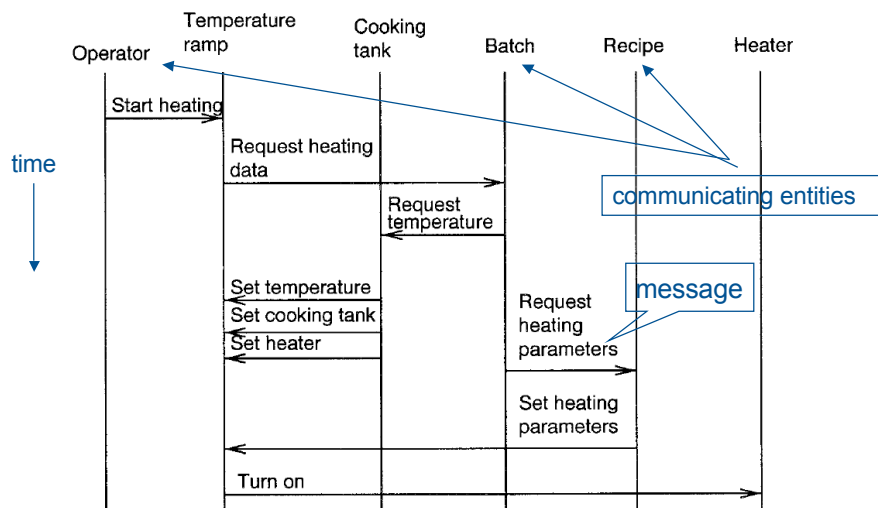
More Diagrams



More diagrammatic notations

- **Object Communication Diagrams**
 - used in the Shlaer–Mellor method to represent object communications
 - nodes represent object classes
 - edges represent possible object communications
 - temporal ordering of communications is not represented.
 - start batch shows a possible communication between an instance of Operator and an instance of Batch
- **JSD System Network Diagrams**
 - used in JSD to specify system functions
 - nodes represent processes
 - directed edges represent communications.
 - processes $\Rightarrow$ **extended finite state machine**, specified by means of a process structure diagram
 - recognize only one kind of component, a process that maintains a state and has behavior over time
 - flat, not hierarchical

Sequence Diagrams

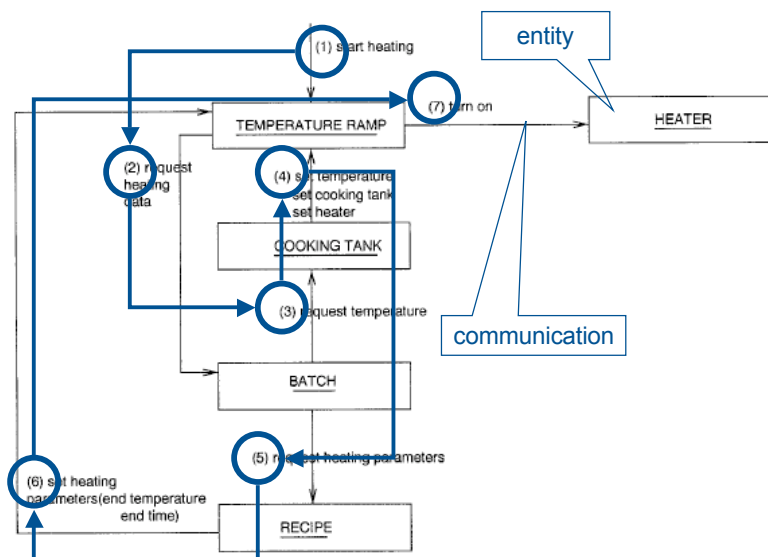


COMPUTER SCIENCE Sequence Diagrams

- represents a particular sequence of messages exchanged between a number of entities
- standardized as message sequence charts in telecom
- popular in object-oriented methods to represent communications between objects
- shows one particular communication sequence in one run of the system
 - show behavior as well as communication
- can be extended with conventions to represent timeouts, global conditions across different entities, delayed message reception, etc.

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COMPUTER SCIENCE Collaboration Diagrams



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Collaboration Diagrams

- popular in object-oriented methods to represent message exchanges between objects
- UML extended with annotations that represent dataflows between the communicating objects and various notations that represent the way in which the communications are implemented.
- differ from other notations
 - nodes represent objects, not activities (DFDs, activity diagrams, activity charts, and block diagrams)
 - nodes represent objects, not object classes (object communication diagrams)
- as in sequence diagrams, represent the sequence of messages in **one particular scenario** whereas all other communication diagrams represent possible communications in **all possible scenarios**.

Representing communication

- two kinds of diagrams to represent communication
 - diagrams that show communication sequence (sequence diagrams and collaboration diagrams)
 - illustrate communications as well as behavior in one particular run of the system
 - diagrams that show a set of possible communications without indicating any sequence (all other diagrams). Diagrams
 - show communication only and do not refer to a particular run of the system.
- the kinds of things that can communicate (conceptual components)
 - .Finite state machines
 - control processes in DFDs are specified by finite state machines.
 - Extended finite state machines
 - contain control, local variables, and tests and updates of these variables.



function specification techniques

- used to specify the external functions of a system
- types
 - Function Refinement Trees
 - Event-Response Specification
 - Use Case Diagrams

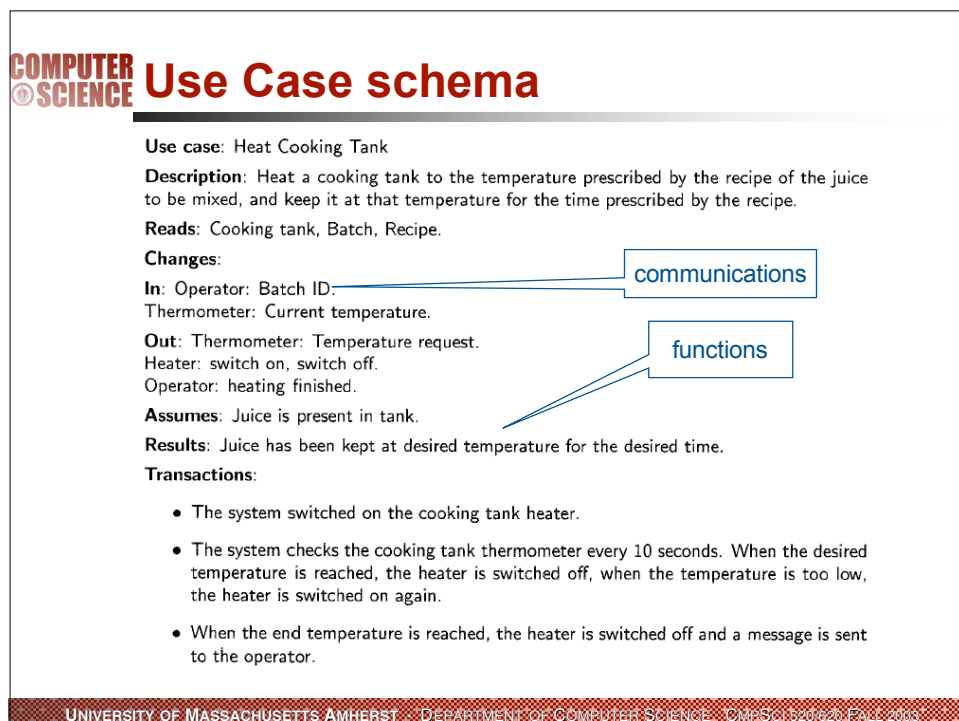
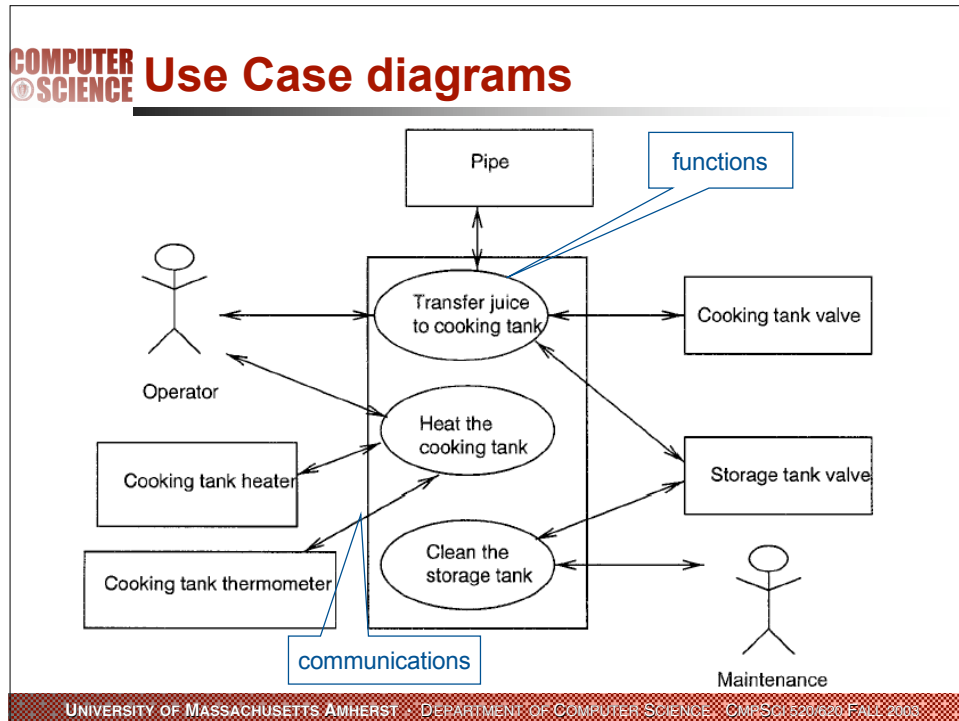
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Use Case diagrams

- represent external system functionality
- a use case
 - an interaction between the system and an external entity that has a use for that external entity
 - need not be atomic
 - can be described by
 - a narrative text
 - by a specification of pre- and post-conditions
 - a state transition diagram.

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function specification techniques

- each technique draws attention to a different aspect of external system functionality
 - no need for formality
 - follow good heuristics
 - don't put more meaning in these diagrams than is intended
- function refinement tree
 - relates the overall mission of a system to its functions down to its atomic transactions.
 - shows why the system must have that function
- Use case diagrams
 - summarize the communications that take place during particular functions.
- context diagram
 - can be used to specify the data exchanges between the system and its environment.
 - draw one partial context diagram for each use case and one for each major system function

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behavior specification techniques

- show how functions of a system or of its components are ordered in time
- Types
 - Process Graphs
 - JSD Process Structure Diagrams
 - Finite State Diagrams
 - Extended Finite State Diagrams
 - Mealy Machines
 - Moore Machines
 - Statecharts
 - SDL State Diagrams
 - Process Dependency Diagrams

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Finite State Machines (FSM's)

- FSM's describe behavior of a system:
 - The sequence of stages/steps/conditions that the system goes through
 - FSM shows how a system acts/reacts to inputs
 - Does this by showing progress through different states
- Hypothesis:
 - The universe in which the system being described must operate can be accurately modeled as always being in exactly one of a finite number of states (situations)
 - There are only a finite number of possible system inputs

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Finite State Machines (FSM)

- finite set of states $S = \{s_1, \dots, s_n\}$
- finite set of inputs $I = \{i_1, \dots, i_n\}$
- transition function $\delta: S \times I \Rightarrow S$
 - can be a partial function
- represent as a graph
 - nodes $\Rightarrow$ states
 - edges $\Rightarrow$ inputs
 - graph $\Rightarrow$ transition function

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Why Use FSM's?

- Primary appeal is visualization
- Intuitively: Can "watch" a stream of inputs "drive" the behavior of the system as a sequence of movements from state to state
- Kinds of questions FSM's seem adept at helping answer:
 - "What is a good way to think about the problem to be solved?"
 - "What is the solution approach?"
 - "How does this program work?"

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Enhancements to FSM's

- Use of hierarchy
- Output annotations on edges
- Distinguished Initial and Terminal states
- Separate data definitions, local and global variables

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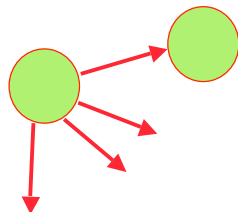
What is FSM good/not good for?

- Focus on specific issue: safety concern
 - Model unsafe state
 - Model state transitions
 - Can unsafe state be reached?
- Drawbacks
 - No sense of functionality
 - No sense of how functionality achieved
 - Difficult and generally impossible to reason about timing

FSM are limited

- Library example
- getbook: index X library $\Rightarrow$ book
- 1,000,000⁺ books in a good library

state = books on shelf

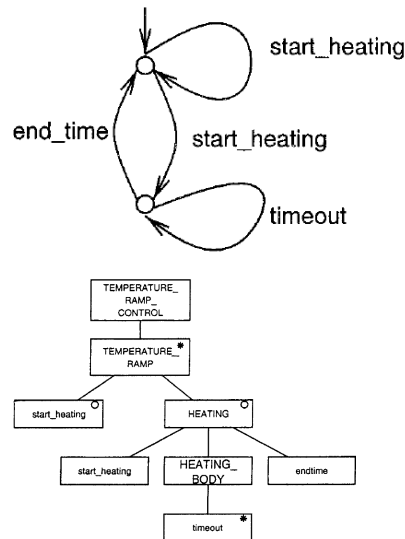


new state = libe - book

2^n transitions could be $\gg 10^6$ states

Some FSM-based notations

- **Process Graphs**
 - nondeterministic
 - may have infinitely many nodes
 - from any node, infinitely many edges may depart
 - used as interpretation structures for formal specifications in process algebra and dynamic logic
- **JSD Process Structure Diagrams.**
 - used in Jackson Structured Programming (JSP)
 - represent the structure of files and of regular programs
 - used in JSD
 - represent the behavior of a system in a modular way
 - a visual way to represent a regular expression by means of a tree diagram



FSM

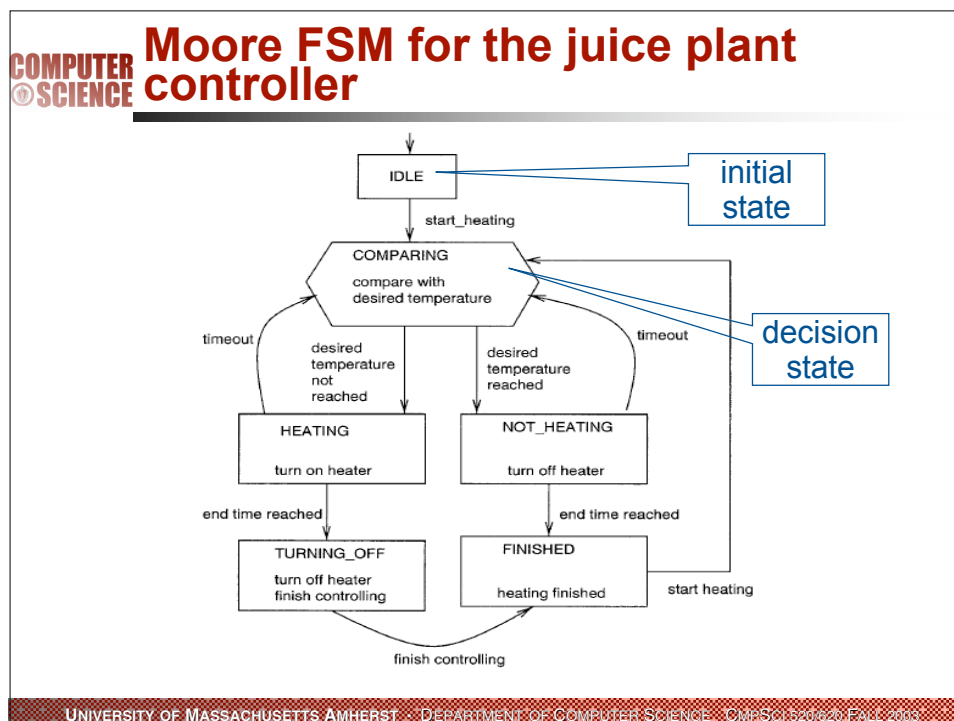
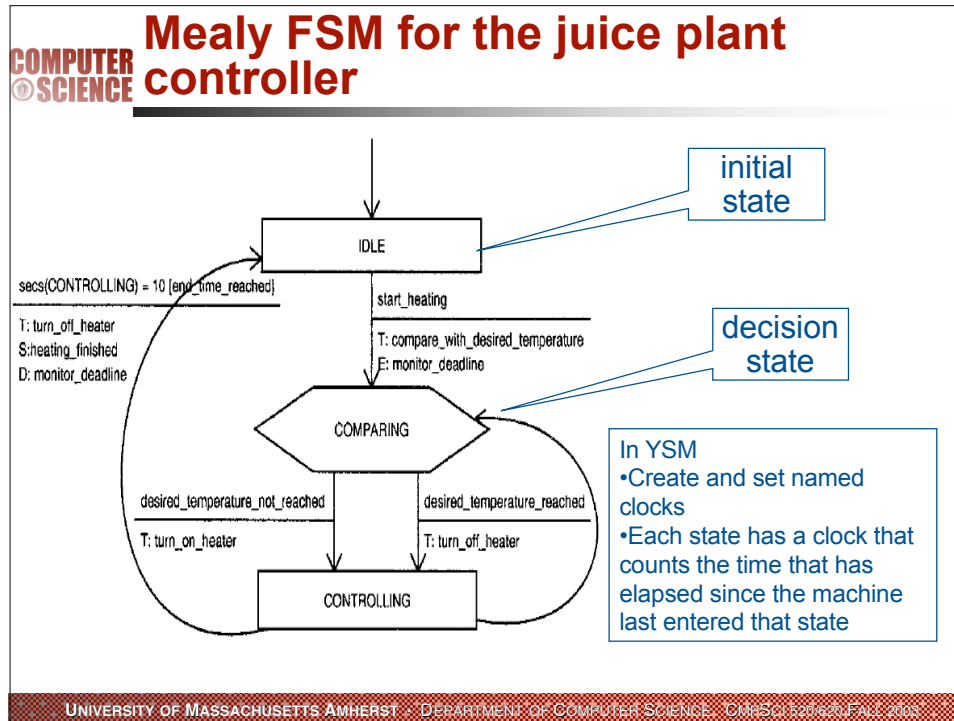
- **Finite State Diagrams**
 - contain only finitely many states and transitions
- **Extended Finite State Diagrams.**
 - number of states can be increased by introducing variables that may be tested and updated by the finite state machine
 - *global state* = *explicit state* (STD nodes) + *extended state* (variables)
 - local variables or external variables
 - local variable is declared together with the specification of the STD + scope rules for these variables (*usually the entire state machine specification*)
 - external variable is declared outside the specification of the STD but can be accessed by means of special operations that act as an interface between the specification and the variables
 - in dataflow models, data stores are external variables with respect to the control processes in the DFD
 - global state change requires communication between the state machine and the external data stores

Extended FSM

- state transition may change the values of variables
- a guard may be specified for each transition that says when the transition can occur.
 - **weak interpretation**,
 - the transition cannot occur if the guard is false
 - guard is in this case a necessary condition of the transition
 - **strong interpretation**,
 - the transition can occur if and only if the guard is true
 - guard is a necessary and sufficient condition of the transition
 - **usually initially specify guards with the weak semantics**
 - when all conditions for a transition are specified, interpret the conjunction of all weak guards as a strong guard
 - a guard could be the conjunction of all preconditions specified for a transition
- include tests in a state machine that are used to determine the next state
 - such tests can be used to resolve nondeterminism
 - a test determines which of a set of possible transitions will occur, thus a test consists of a guard for each of the possible transitions.

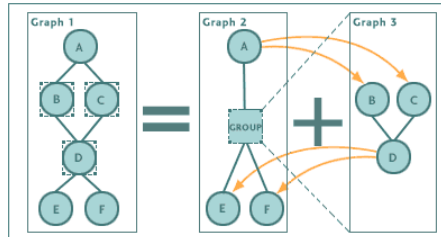
Mealy & Moore Machines

- Mealy machine
 - output actions are associated with transitions
- Moore Machines
 - outputs are associated with states



COMPUTER SCIENCE Statecharts

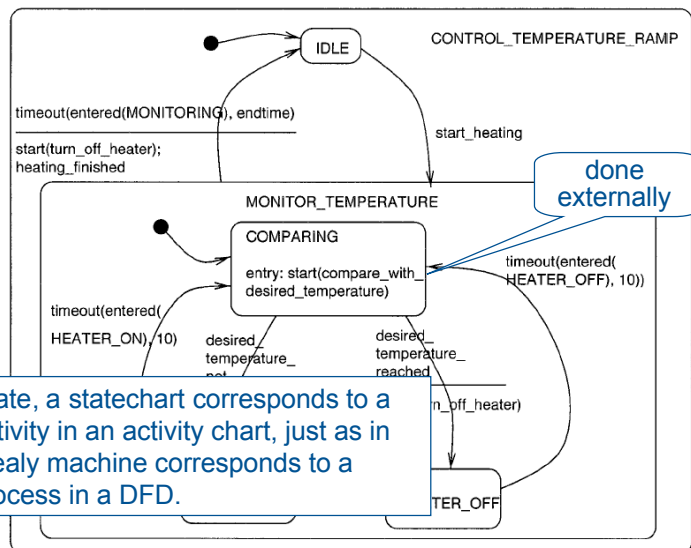
- higraph without intersection but with Cartesian products



- node inclusion allows us to partition a state into substates
- Cartesian products allow us to specify parallelism
- actions can be specified
 - along transitions (Mealy)
 - upon entry of states (Moore), and
 - exit of states
- local variables represent the extended state.

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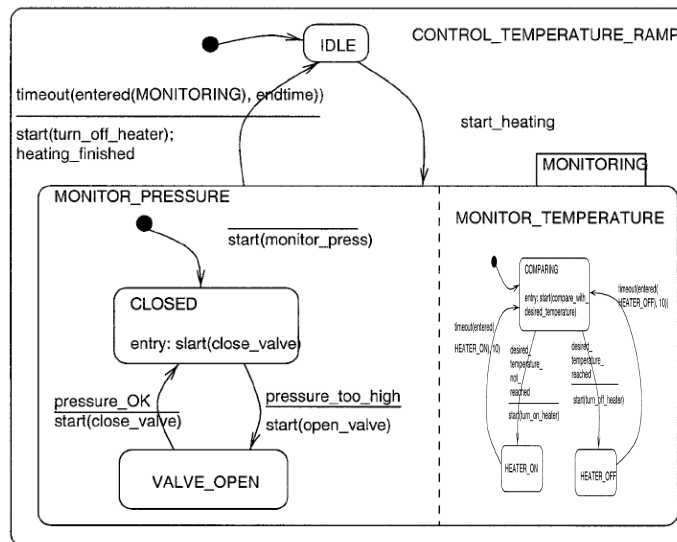
COMPUTER SCIENCE Statecharts



In Statestate, a statechart corresponds to a control activity in an activity chart, just as in YSM a Mealy machine corresponds to a control process in a DFD.

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Statechart with parallelism

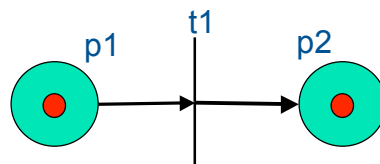


FSMs

- different notations are very similar
 - all represent states by labeled nodes and transitions by labeled directed (hyper)edges
 - all except process dependency diagrams have a formal semantics
 - many different possible semantics of statecharts and statechart-like notations
- time in state machines.
 - point or interval semantics of time
 - discrete, dense, or continuous model of time, and assume that transitions take time or are instantaneous
 - these choices leads to important differences in the behavior specified by a state diagram
 - semantics of time is carefully defined in SDL and in the Statemate semantics of statecharts.

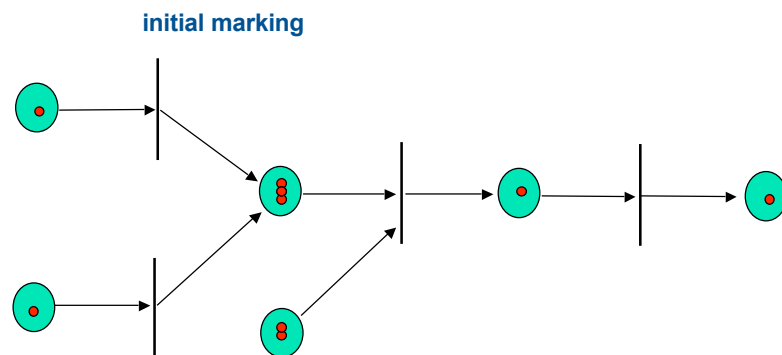
Petri Nets

- Petri nets are “marked” graphs
 - two node types: places & transitions
 - tokens mark the nodes
 - transitions are enabled (“fire”) if all connected places contain tokens



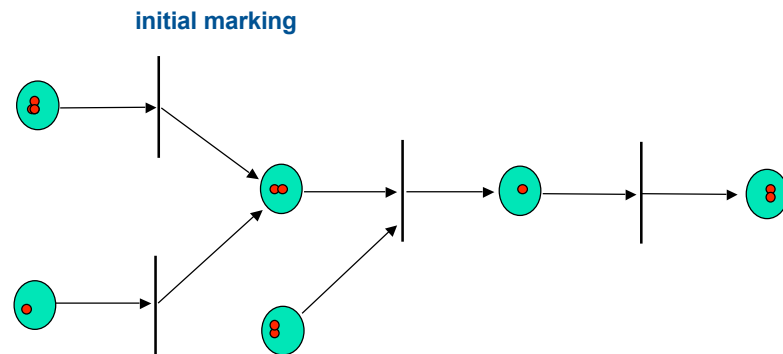
Petri Nets

- Example - simultaneous firing



COMPUTER SCIENCE Petri Nets

- Example - asynchronous, atomic firing



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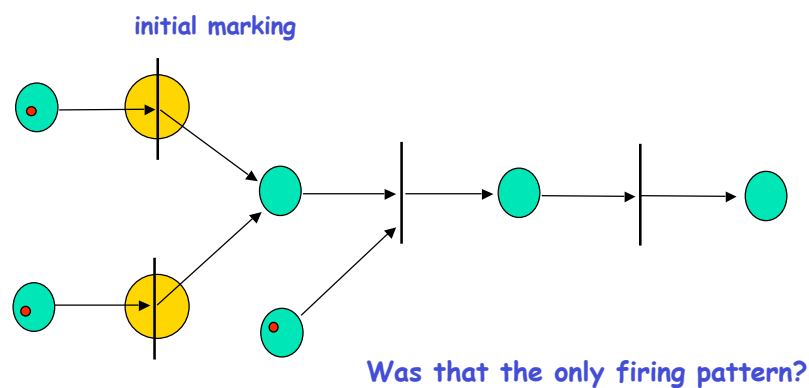
COMPUTER SCIENCE Petri Nets: Informal Definition

- Designed specifically for modeling systems with interacting concurrent components.
- Consists of a set of places and a set of transitions
 - Edges connect places and transitions.
 - Only transition $\rightarrow$ place and place $\rightarrow$ transition links are allowed.
- Each place can have a finite number of tokens.
- A transition is enabled if each of its input places has at least one token.
 - An enabled transition can fire: one token is taken from each input place and one token is put into each output place.

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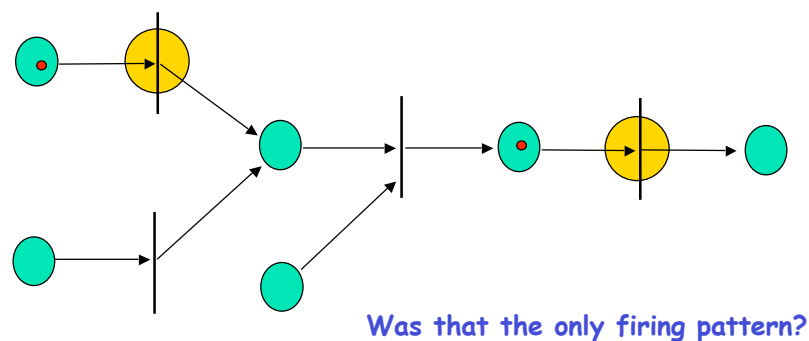
Petri Nets

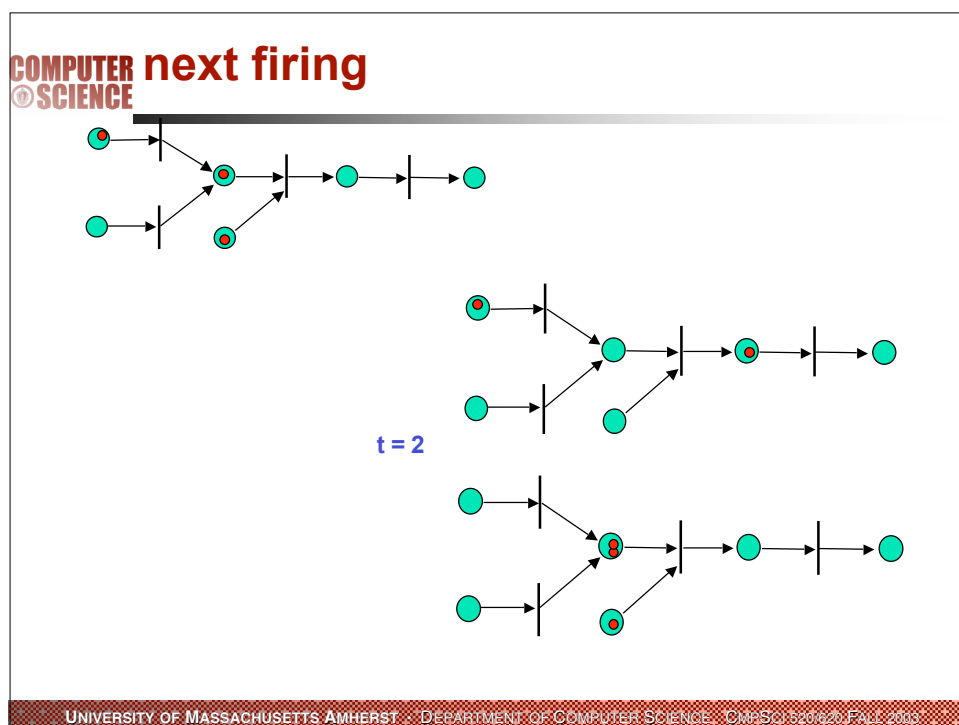
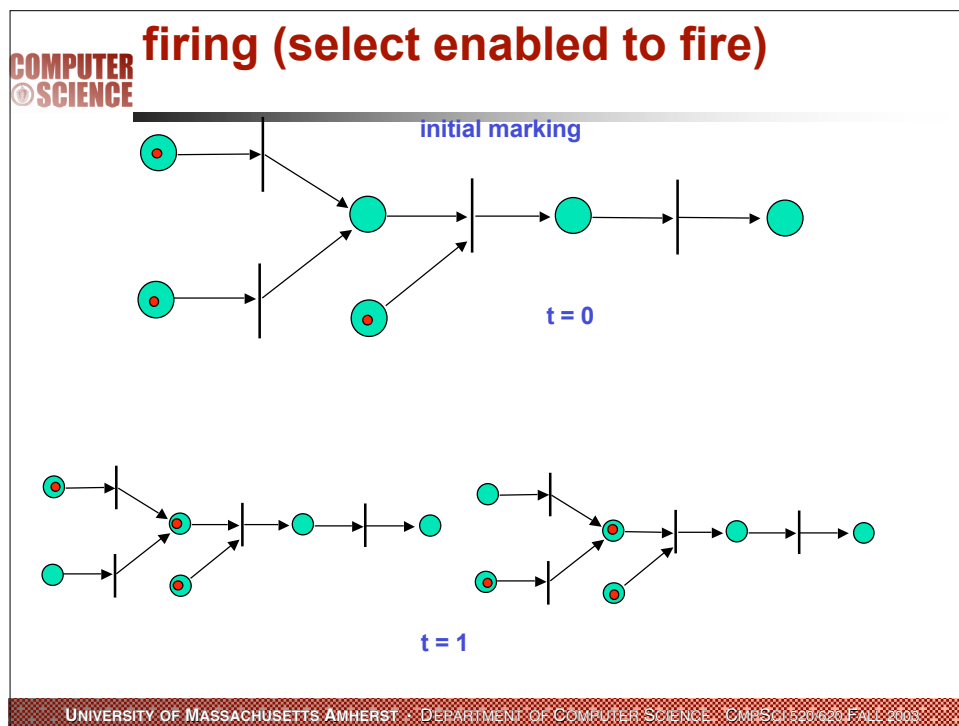
- Example - asynchronous, atomic firing

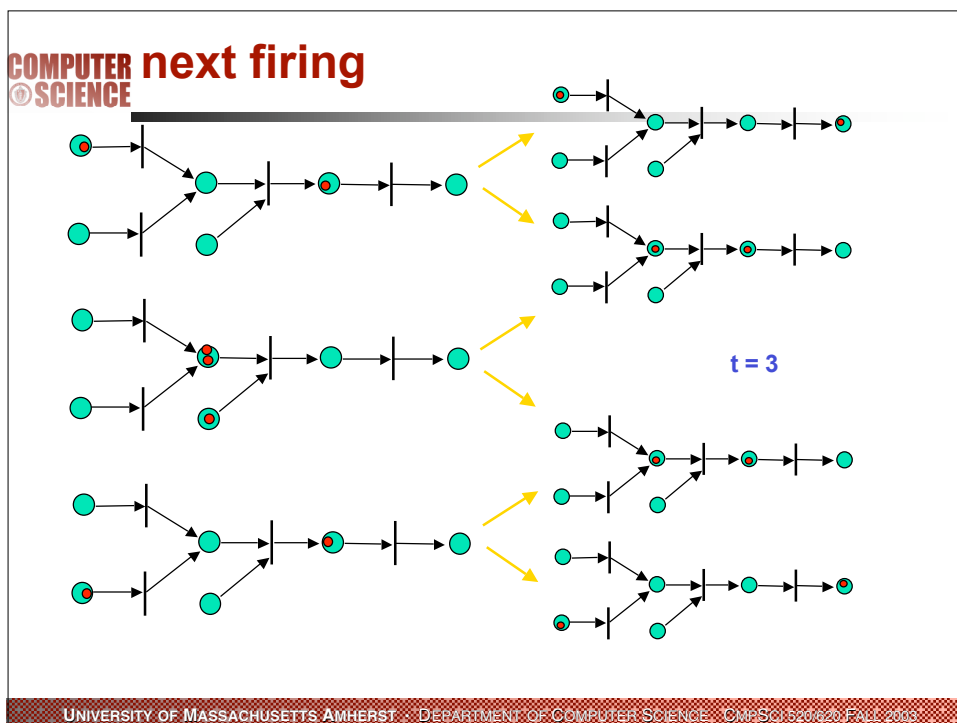
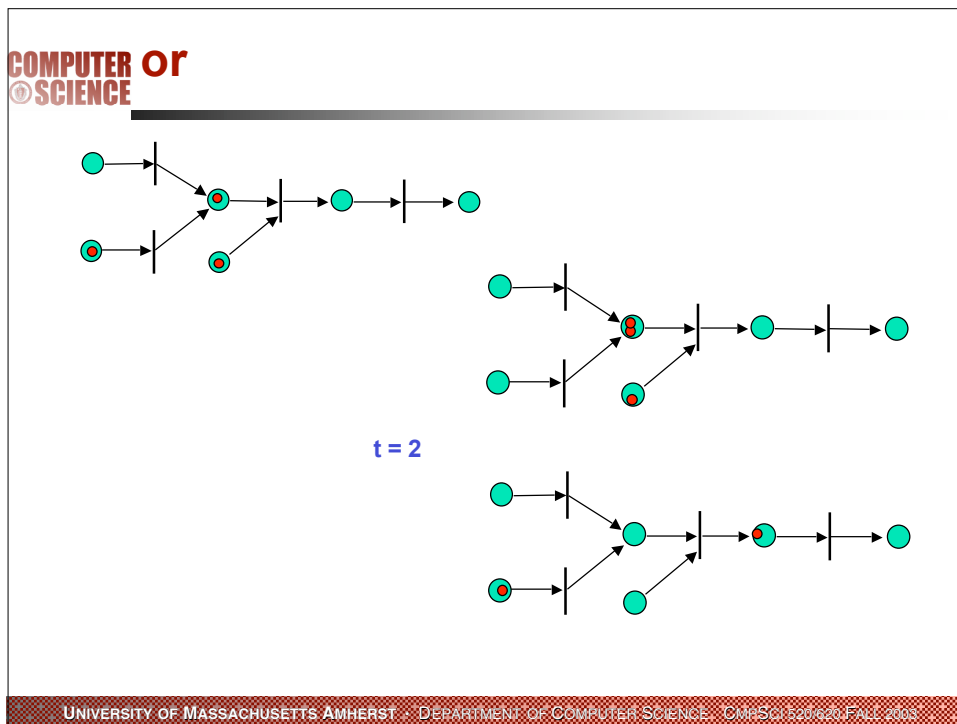


Petri Nets

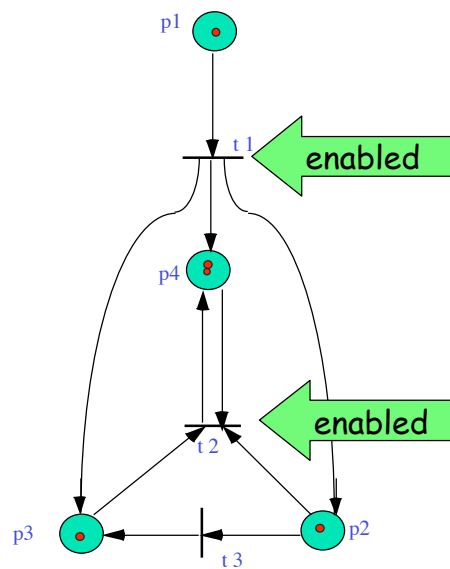
- Example - asynchronous, atomic firing





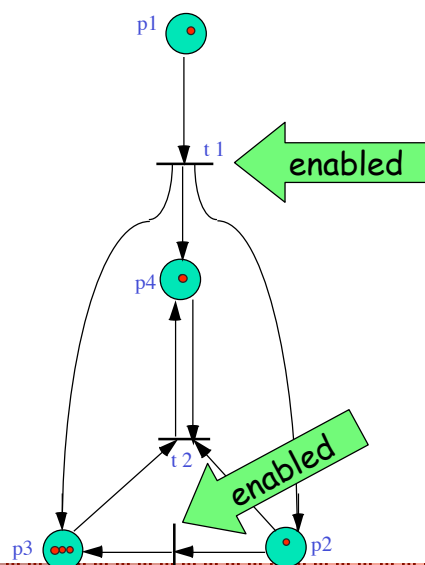


Petri Net example



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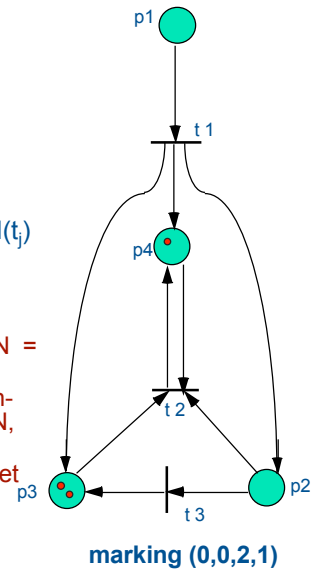
Petri Net example



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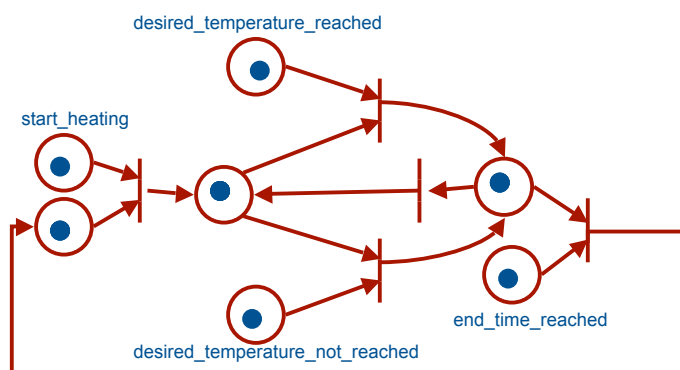
Petri Nets: Formal Definition

- A Petri Net is a four-tuple, $C=(P,T,I,O)$
- $P = \{p_1, p_2, \dots, p_n\}$, $n \geq 0$ is a finite set of places.
- $T = \{t_1, t_2, \dots, t_m\}$, $m \geq 0$ is a finite set of transitions.
- $I: T \rightarrow P$ is the input function.
- $O: T \rightarrow P$ is the output function.
- p_i is an input place of a transition t_j if $p_i \in I(t_j)$
- p_i is an output place of a transition t_j if $p_i \in O(t_j)$
- Petri Net markings
 - A marking m is a mapping $P \rightarrow N$ where $N = 0, 1, 2, \dots$
 - The marking m can be represented as a n -vector $m = (m_1, m_2, \dots, m_n)$, $n = |P|$, $m_i \in N$, $1 \leq i \leq n$
 - A marked Petri net $M = (C, m)$ is a Petri net C and a marking m .



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Petri Net for Heating Controller



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More on Petri nets

- if there exists a marking which is reachable from the initial marking where no transitions are enabled, such a transition is called a "deadlock"
- a PN with no possible deadlock is said to be live, called the "liveness property"
- in simplest PN, tokens are uninterpreted
 - in general, a selection policy can not be specified
 - have no "policy" for resolving conflicts, potential "starvation"
- many extensions:
 - Hierarchical Petri Nets
 - Colored tokens
 - "Or" transitions
 - Queues at places

Petri Nets vs. Finite State Machines

- For any finite state machine, a Petri net can be built that models the finite state machine
 - Petri nets are as powerful as finite state machines
- Petri nets advantages:
 - net composition (in different forms) can be found easier than the cross-product of finite state machines
 - parallelism and nondeterminism are represented in a more understandable way
- Finite state machines advantage:
 - simpler graph structure for some applications (e.g. parsers)